Hyper-Minimization in $O(n^2)$

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What is hyper-minimization?

Classical Equivalence

- D₁ and D₂ recognize the same language.
- $L(D_1) \otimes L(D_2)$ is empty
- Notation: $D_1 \approx D_2$

What is hyper-minimization?

| Classical Equivalence | F-Equivalence |
|--|---|
| D₁ and D₂ recognize the same language. | D₁ and D₂ "almost" recognize the same language. |
| • $L(D_1) \otimes L(D_2)$ is empty | • $L(D_1) \otimes L(D_2)$ is finite |
| • Notation: $D_1 \approx D_2$ | • Notation: $D_1 \sim D_2$ |
| | |
| | |

Small Example







Showing Myhill-Nerode equivalence classes





Showing f-equivalence classes





Showing Myhill-Nerode equivalence classes

Small Example – classically minimized



Small Example – classically minimized



Showing f-equivalence classes

Small Example – hyper-minimized



Showing f-equivalence classes

Small Example – hyper-minimized



Elementary properties

• Let q_1 be a state from DFA D_1 , and q_2 be a state from D_2 . If $q_1 \sim q_2$, then for any input $c, \delta(q_1, c) \sim \delta(q_2, c)$.

• If $D_1 \sim D_2$, then $\forall q_1 \in Q_1$, $\exists q_2 \in Q_2$: $q_1 \sim q_2$.



Kernel isomorphism

If $D_1 \sim D_2$ and both are classically minimized, then their kernels are isomorphic.



Preamble isomorphism

If $D_1 \sim D_2$ and both are hyper-minimized, then their preambles are (somewhat) isomorphic. These aspects within the preamble may differ:

- Whether a preamble state is accepting or not.
- Transitions from the preamble to the kernel can move within an f-equivalence class.

Minimization Algorithm

Classical Minimization

- I. Delete unreachable states
- 2. Find equivalent states
- 3. Merge states within each equivalence class

Minimization Algorithm

| Classical Minimization | Hyper-Minimization |
|---|---|
| I. Delete unreachable states | I. (Classically) Minimize |
| 2. Find equivalent states | 2. Find equivalent states |
| 3. Merge states within each equivalence class | 3. Merge states within each equivalence class |

I. Classically minimal



2. Finding f-equivalent state classes

- i. Let $D_{\otimes} = D \otimes D$ be the standard DFA crossproduct construction for symmetric difference.
- ii. Find all states (q_0, q_1) in D_{\otimes} which induce a finite language q_0 and q_1 are f-equivalent in D.
- iii. Use the list of these pairs to construct the equivalence classes.

2.i. Cross-product with self



2.ii. Find all right-finite states



2.ii. Find all right-finite states



2.iii. Construct equivalence classes from pairs



2.iii. Construct equivalence classes from pairs



Showing f-equivalence classes

3. Merge states within each equivalence class



Showing f-equivalence classes

Finite-Factoring

- Motivation: hyper-minimization changes the language
- Use hyper-minimization to split a regular language into two parts: infinite and finite
- Use a DFCA to recognize the finite part
- What is a DFCA?

Finite-Factoring Algorithm

- I. Let $D' = hyper_minimize(D)$
- 2. Let $D_f = xor_cross_product(D, D')$
- 3. Let $n = max(|w| : w \in L(D_f))$
- 4. Minimize the DFCA (D_f, n)
- 5. Return $(D', (D_f, n))$

Finite-Factoring



Finite-Factoring





Open Problems, Questions

Source code: <u>http://ianab.com/hyper/</u>

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